POSTER: Optimizing Sparse Computations Jointly

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Abstract

This work proposes a framework called FuSy that analyzes the data dependence graphs (DAGs) of two sparse kernels and creates an efficient schedule to execute the kernels in combination. Sparse kernels are frequently used in scientific codes and in machine learning algorithms and very often they are used in combination. Iterative linear system solvers are an example where kernels such as sparse triangular solver (SpTRSV) and sparse matrix-vector multiplication (SpMV) are called consecutively in each iteration of the solver. Prior approaches typically optimize these sparse kernels independently leading to high synchronization overheads and low locality. We propose an approach that analyzes the DAGs of two sparse kernels and then creates a new order of execution that enables running the two kernels efficiently in parallel. To investigate the efficiency of our approach, we compare it with the state-of-the-art MKL library for two kernel combinations, SpTRSV-SpMV and SpMV-SpTRSV which are commonly used in iterative solvers. Experimental results show that our approach is on average $2.6 \times$ and $1.8 \times$ faster than the MKL library for a set of matrices from the Suitesparse matrix repository.

CCS Concepts: • Computing methodologies \rightarrow Shared memory algorithms; • Software and its engineering \rightarrow Source code generation.

Keywords: Sparse matrix code, loop fusion, loop-carried dependence

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1 Extended Abstract

The numerical methods [10] that are frequently used in realworld applications such as in scientific simulations and data analytic codes are often composed of a number of sparse matrix computations that execute inside an iteration of the numerical algorithm and between iterations. Because sparse kernels are often the most time-consuming operation in these applications, numerous library and compiler approaches have been proposed to optimize these kernels. However, prior work primarily optimizes sparse kernels in isolation thus when used to accelerate real-world simulations the realized speedups are sometimes not significant.

Numerous parallel sparse libraries [11, 16] and inspectorexecutor approaches that inspect memory access patterns at runtime such as [6, 8, 12, 15] optimize individual sparse matrix kernels. Kernels with a fully parallel outermost loop have sufficient parallelism and thus an efficient schedule is needed to create a balanced parallel implementation. Sparse kernels with partial parallelism, i.e. loop-carried dependencies, have irregular computation patterns that depend on the sparse matrix code and input data, thus runtime inspection is required to extract the computation patterns. In inspectorexecutor frameworks and libraries such as [2, 13], a data flow directed acyclic graph (DAG) is built to expose data dependencies. For example, the inspectors in [6, 15], use wavefront parallelism to create a parallel schedule for kernels with partial parallelism. First, the DAG is created, and is traversed in topological order to create a list of wavefronts that are iterations that can execute in parallel; this is known as wavefront parallelism.

Wavefront parallelism requires synchronization between wavefronts, and thus when applied to individual sparse kernels with loop-carried dependencies can be less efficient due to synchronization overheads. Also, for sparse kernels with non-uniform workloads, such as Cholesky [3], wavefront methods can lead to load imbalance.

DAG partitioning techniques such as DAGP [7] (used in [9]) typically create fewer wavefronts, thus reducing synchronization overheads, and group iterations that reuse data to improve data locality. DAGP adopts a multilevel approach [1] with coarsening and refinement for acyclic partitioning of DAG. These techniques are efficient for individual sparse computations, however, when applied to the joint DAG, they create some non-linear overheads or large exploration space and thus significantly increase analysis time. For example,

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when applied to the joint DAG of two sparse kernels such as sparse triangular solver (SpTRSV) and sparse matrix-vector multiplication (SpMV), DAGP becomes 5 times slower for selected set of matrices from Suitesparse [5], even though the joint DAG size increases two times.

We present an approach that creates an efficient schedule for when sparse kernels are used jointly. Our approach analyzes the data dependence graph of two sparse kernels to create a load-balanced parallel schedule of the combined code with good locality. Locality is improved by assigning dependent vertices to the same group and then executing that group of vertices via the same thread. Load balance is improved by assigning vertices throughout execution to create well-balanced tasks.

Results are collected on a Haswell multicore architecture with 12 cores of a Xeon E5-2680v3 processor and a 30MB L3 cache. We use matrices from the Suitesparse [5] repository to compare our approach with MKL [16]. The matrices are selected to be of different sizes and varying sparsity patterns from a small number of nonzero elements (1.4×10^5) to a large number of nonzero elements (1.1×10^8) . We use MKL 2019.3.199 and call each kernel from the MKL library separately. The performance of both our approach and MKL are tested on two kernel combinations, SpTRSV-SpMV and SpMV-SpTRSV. Both combinations are used in iterative linear solver methods such as preconditioned GMRES [4] and Gauss-seidel. Here we focus on joint optimization of kernels within an iteration of the solver to ensure the stability of the solver. For the combinations of SpTRSV-SpMV and SpMV-SpTRSV, we improved over the MKL library by an average speedup of $2.6 \times$ and $1.8 \times$ consecutively.

This work focuses primarily on optimizing the combination of SpTRSV and SpMV using the proposed approach. As future work, the DAGs of other sparse kernels and their combination should be studied to create an efficient schedule for joining these kernels. Also, a combination of some sparse kernels might not be cost-efficient, and hence new models should be investigated for the efficiency of joint execution. We also plan to apply the proposed approach to real-world benchmarks to demonstrate the effect of sparse kernel combinations.

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